

## FIFOs

- The specification of <u>mode argument</u> is the <u>same as for the open system call</u>.
- The rule for user and group ownership of a FIFO are the same as in a file.
- Once we create a FIFO by using mkfifo, we can open it by using open().

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## FIFOs

- A FIFO supports blocked read and write operations by default: if a process opens the FIFO for reading, it is blocked until another process opens the FIFO for writing, and vice versa.
   However, it is possible to make FIFOs support
- However, it is possible to make FIFOs support non-blocking operations by specifying the O\_NONBLOCK flag while opening them.
- A FIFO must be opened either read-only or writeonly. It must not be opened for read-write because it is half-duplex, that is, a one-way channel.

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# FIFOs

- In the normal case (O\_NONBLOCK not specified)
  - <u>An FIFO</u> open for read-only <u>blocks</u> until some other process opens the FIFO for writing.
  - <u>An FIFO</u> open for write-only <u>blocks</u> until some other process open the FIFO for reading.
- If O\_NONBLOCK is specified
  - An open for read-only returns immediately.
  - An open for write-only returns –I with errno set to ENXIO if no process has the FIFO open for reading.

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# FIFOs

- FIFOs are used by shell commands to pass data from one shell pipleline to another without creating intermediate temporary files.
- FIFOs are used as rendezvous point in client-server applications to pass data between the clients and the servers.









ce(1, inverse argumence(names two integers(h", 37);	
te (1 "involid commente) eRetex tue internals" (27).	
e(1, line, strien(line));	
ntr(line, "The sum is ed(h", intivint2);	
and other and the second state of the second state of the	
e, "ed ed", fint1, fint2)== 2)	
line, BFFERSIZE))>0)	
8 turn p	
EX, O_RDONLY);	
r reading */	
Felo //	
reating the named nine").	
45 (errec 1= REVIST)) /	
LF_DUFLEX, U666);	
- pipe */	
12	
unt, numread;	
*arov[]]	
20	
"halfduplex"	
	"haldhglas" 28 29 *apqull ant, unseed; ; ; pips */ g_DURK_066;; at (seroe i= EDITY) ( teating the named pips"; r scaling */ g, 0_LORD(1); ( teating the named pips"; r scaling */ g, 0_LORD(1); s,



# XSI Interprocess Communication

- There are three types of XSI IPC
  - Message queue
  - Semaphore
  - Shared memory
- Each IPC structure in the kernel is <u>referred to</u> by a **non-negative identifier**.
- When a given IPC structure is created and then removed, the identifier associated with that structure continually increase up to the maximum positive integer, and then wraps around to 0.

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# XSI Interprocess Communication

- When an XSI IPC structure is created (by calling msgget(), semget() or shmget()), a key must be specified.
- The data type **key\_t** for a key is specified in the header file <sys/types.h>.

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# XSI Interprocess Communication

- The ftok() function return a key based on path and id that is usable in subsequent calls to msgget(), semget(), and shmget().
- The application shall ensure that the *path* argument is <u>the pathname of an existing file</u>.
- □ Only lower 8 bit of id are used when generating a queue (∴we can path a character).

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# XSI Interprocess Communication

**•** XSI IPC <u>associated with ipc perm structure</u>.

□ This structure defines the permissions and owner and so on.

# struct ipc\_perm { uid\_t uid; /\* owner's effective user ID\*/ gid\_t gid; /\* owner's effective group ID \*/ uid\_t cuid; /\* creator effective user ID \*/ gid\_t cgid; /\* creator effective group ID \*/ mode\_t mode /\* access mode \*/ };

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### XSI Interprocess Communication All the fields are initialized when the IPC structure is created. We can modify the uid, gid, and mode filed by calling msgctl(), semctl() or shmctl(). The value of mode fields are: Operation permissions Octal value Read by user 00400 = 000.000 100.000.000 Write by user 00200 = 000 000 010 000 000 Read by group 00040 = 000 000 000 100 000 Write by group 00020 = 000 000 000 010 000 00004 = 000 000 000 000 100 Read by others $00002 = 000\,000\,000\,000\,010$ Write by others













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### XSI IPC (Message Queue) XSI IPC (Message Queue) #include <sys/msg.h> int msgrcv(int msqid, void \*msgp, int msgsz, long msgtyp, int msgflg); Return key if 0k, -1 error The msgrcv() function reads a message from the queue msgtyp Specifies the type of message requested as associated with msqid and places it in the user-defined follows: structure that *msgp* points to. <u>If msgtyp is 0</u>, the first message on the queue is received. • If msgtyp is greater than 0, the first message of type equal to msgtyp is received. When successfully completed, the following actions are taken with respect to the data structure associated with If msgtyp is less than 0, the first message of the lowest type that is less than or equal to the absolute value of msgtyp is received. msqid: msg\_qnum is decremented by 1. msaflq Specifies the action to be taken if a message of the msg\_lrpid is set to the process ID of the calling process. desired type is not in the queue. msg\_rtime is set to the current time. COSC350 System Software, Fall 2024 Dr. Sang-Eon Park 25 COSC350 System Software, Fall 2024 Dr. Sang-Eon Park





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	/*spock.c read message from the message queue */	
	/* created by kirt.c */	
	<pre>#include <stdio.h></stdio.h></pre>	
	<pre>#include <stdlib.h></stdlib.h></pre>	
	<pre>#include <errno.h></errno.h></pre>	
	<pre>#include <sys types.h=""></sys></pre>	
	<pre>#include <sys ipc.h=""></sys></pre>	
	<pre>#include <sys msg.h=""></sys></pre>	
	struct my_msgbuf (	
	long mtype;	
	char mtext[200];	
	12	
	int main (vold)	
	struct my magdur bur)	
	Int maging	
	velt velt	
	/* create a key same as kirt c */	
	if $(lev = frok("kirk c", 'B')) = -1)$	
	nerror ("ftok erne")	
	evit(1);	
	/* open message queue already created by kirk.c */	
	if ((msqid = msqqet(key, 0644)) == -1) [	
	perror("msqqet error");	
	exit(1);	
	3	
	printf("spock: ready to receive messages, captain.\n");	
	for(11) (	
	/* get each message from the message queue */	
	if (msgrcv(msqid, (struct msgbuf *)&buf, sizeof(buf), 0, 0) == -1) {	
	perior("msgrcv error");	
	exit(1);	
	printr("spock: \"%s\"\n", bur.mtext);	
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# XSI IPC (Shared Memory)

- Shared memory can be faster than message passing, since message-passing systems are typically implemented using system calls (shared memory are located in user's space).
- In shared-memory systems, system calls are required only to establish shared memory regions.
- Once shared memory is established, all accesses are treated as routine memory accesses, without kernel's assistance.

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# XSI IPC (Shared Memory)

- Shared memory allows two or more processes to share a given region of memory. This is the fastest form of IPC, because the data does not need to be copied between the client and the server (or between processes).
- The only trick in using shared memory is synchronizing access to a given region among multiple processes.
- Since OS does not support mutual exclusion, programmer must take care mutual exclusion of the region between multiple processes by using a semaphore.
- $\blacksquare$  The kernel maintains a structure <code>shmid\_ds</code> with at least the following members for each shared memory segment:

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# XSI IPC (Shared Memory) Before using the shared memory what we needs to be done with the system calls, Create the shared memory segment or use an already created shared memory segment (shmget()) Attach the process to the already created shared memory segment (shmat()) Detach the process from the already attached shared memory segment (shmdt()). Control operations on the shared memory segment (shmctl())



### XSI IPC (Shared Memory) XSI IPC (Shared Memory) Once a shared memory segment has been created, a process attaches it to its address space by calling system call $\tt shmat()$ . #include <sys/shm.h> #include <sys/sum.... #include <sys/ipc.h> int shmget (key\_t key, size\_t size, int shmflg); Return shared memory ID if Ok, -1 error #include <sys/ipc.h> void \*shmat(int shmid, const void \*addr, int flag); □ If a new shared memory is created, the ipc perm Return the address of attached shared memory if Ok, -1 error structure are initiated. shmid: ID return by shmget() system call. shm\_lpid, shm\_nattch, shm\_atime, and shm\_dtime are all set to 0. addr: is to specify the attaching address. If addr is shm\_ctime is set to the current time. NULL, the system chooses the suitable address to attach the segment by default. If it is not NULL and SHM RND is shm\_segsz is set to the size requested. specified in flag, attach is equal to the address of the nearest multiple of SHMLBA(Lower Boundary Address). COSC350 System Software, Fall 2024 Dr. Sang-Eon Park 37 COSC350 System Software, Fall 2024 Dr. Sang-Eon Park

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XSI IPC (Shared Memory)		
// header.h		
define NOT	READY -1	
define FILJ	LED 0 //when sender fill data	
define TAKJ	EN 1 //when receiver take data	
#define GO	2 // when sender keep sending	
#define STOI	P 3 // when sender stop sending data	
struct studer int char char	nt ( id; : lname[20]; : fname[20];	
struct Memory	V {	
int sta	atus; //FILLED or TAKEN	
int gor	stop; //GO or STOP	
struct #	student data;	
};		
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}	Tettin o,	
	shm->gostop = GO;	
	shm->status = NOT_READY;	
	}	
	exit (1);	
	perror("shmat error \n");	
	11 ((10ng) Snm == -1)	
	<pre>shm = (struct Memory *) shmat(shmid, NULL, 0); //attach to shared memory if (() and box</pre>	
	}	
	exit (1);	
	<pre>perror("shmget error \n");</pre>	
	{	
	<pre>//create a suareu memory if ((ehmid = ehmmat(kau, eizenf(etrunt Memoru) IDC CDEAT   0666)) (0)</pre>	
	<pre>key = ftok(".", 'x'); //create a key value //oreate a shared momentum</pre>	
	struct Memory *shm;	
	key_t key;	
	int shmid;	
{		
int mai	(int argc, char *argv[])	
#includ	e "neader.n"	
finclud	e <errno.h></errno.h>	
<pre>#include<sys shm.h=""></sys></pre>		
<pre>#include<stdlib.h></stdlib.h></pre>		
#includ	e <stdio.h></stdio.h>	
//build	em eh	











int shmid;
int shmidy
Any_b Any, abar Tahn Tah
Long - Stady - Sy Long - Stady H - Hally //assats a long value
//reate a shared merry with size 100 byte
if ((sheid = sheart(she) sheart(street, tor (sheart   0.666)) c0)
(aming - amperiary, and m, includer ( obody) (a)
t measure (Polyment, ensure ) off) -
parts ().
ware (i),
shw = shwat(shwid, NULL, 0) - //attach pointer to the shared memory
if (share (shart) =1)
ar (sum - (cont ) -1)
nonconfidence on the later
parts ().
WALC (1)
<pre>/ memopy (shm, "Mello World", 11); //write to shared memory you can use write system call</pre>
s = shm;
s+=11;
*8 = 0;
<pre>while (*shm != '*') //server will wait until client read and type * in shared memory</pre>
printf("Server has detected the completion of its child\n");
shedt((vold ') she); //detach shared menory
print! ("Server has detached its shared memory\n");
shmcti(shmid, IPC_RHID, NULL); //remove shared memory
printf("Server has removed its shared memory\n");
return 0;

// client	i.c ; open shared memory and read data
#include+	<pre>(string.h&gt;)</pre>
#include <sys ipc.b=""></sys>	
+includes/systems b>	
fincture(sys) cypester	
#define 3	BHSIZE 100
int main	(int argc, char *argv[])
	int shmid;
	key t keyj
	char *shin, *s;
	<pre>key = ftok(".", 'x'); //create a key value</pre>
	if ((shmid = shmget(key, SHSIZE,0666)) <0) //open shared variable created by server
	perror('sniget error (n'))
	shm = shmat(shmid, NULL, 0); // attach a pointer to shared memory
	if (shm == (char*) -1)
	6
	perror("shmat error \n");
	exit (1);
	for (s =shm; *s != 0; s++) //read available data from the shared memory
	<pre>printf("%c", *s);</pre>
	printf("\n");
	*shm = '*'; // write a `*' to shared memory which inform to server that client done its job
	return 0;
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# XSI IPC (Semaphore) To access a shared resources, a process needs to do the followings: Test the semaphore that controls the resources. If the value of the semaphore is >0, the process reduce the value by 1 and access resources. Check and modification to the a semaphore are executed indivisibly. If the value of the semaphore is 0, the process need go to sleep on the semaphore until the value becomes greater than 0.

# XSI IPC (Semaphore) Ex) Lets there are two processes P<sub>1</sub>, P<sub>2</sub> working on their job and , and two resource R<sub>1</sub> and R<sub>2</sub>. Both P<sub>1</sub> and P<sub>2</sub> need R<sub>1</sub> and R<sub>2</sub> to finish their job. Each resource is associated with a semaphore.

XSI IPC (Semaphore)	XSI IPC (Semaphore)
Case 1) semaphore R <sub>1</sub> ; semaphore R <sub>2</sub> ; void process_P <sub>1</sub> () void process_P <sub>2</sub> () { down(&R <sub>1</sub> ); down(&R <sub>2</sub> ); down(&R <sub>2</sub> ); use_both_resource(); use_both_resource(); up(&R <sub>2</sub> ); up(&R <sub>2</sub> ); up(&R <sub>1</sub> ); up(&R <sub>1</sub> ); } }	Case 2) semaphore R <sub>1</sub> ; semaphore R <sub>2</sub> ; void process_P <sub>1</sub> () void process_P <sub>2</sub> () {
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