## COSC 220: Computer Science II

## Module 1

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Spring 2021

## Arrays

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### 1.1 Array in C++

- Array: variable that can store multiple values of the same type
- Values are stored in adjacent memory locations
- Declared using [] operator:



## Accessing Array Elements

- Each element in an array is assigned a unique subscript from 0 to $n-1$
- Access an element in an array:

```
array_name[subscript]
```

int tests[5];


## Accessing Array Elements (Cont'd)

- Each array element can be used as a regular variable:

```
tests[0] = 79;
cout << tests[0];
cin >> tests[1];
tests[4] = tests[0] + tests[1];
```

- Arrays must be accessed via individual elements:
cout << tests; // not legal


## Using a Loop to Step Through an Array

- Example - The following code defines an array, numbers, and assigns 99 to each element:
const int ARRAY_SIZE = 5;
int numbers[ARRAY_SIZE];

$$
\begin{gathered}
\text { for (int count }=0 ; \text { count < ARRAY_SIZE; count++) } \\
\text { numbers [count] }=99 ;
\end{gathered}
$$

The variable count starts at 0 , which is the first valid subscript value

The loop ends when the variable count reaches 5 , which is the first invalid subscript value

The variable count is incremented after each iteration

## Array Initialization

- An array can be initialized with an initialization list:
const int SIZE = 5;
int tests[SIZE] $=\{79,82,91,77,84\}$;
- The values are stored in the array in the order in which they appear in the list.

- The initialization list cannot exceed the array size.


## No Bounds Checking in C++

- When you use an array subscript, C++ does not check whether it is a valid subscript or not
> You can use subscripts that are beyond the bounds of the array
int values [3] $=\{5,8,10\}$;
// Syntax correct, but may corrupt other memory
// locations, crash program, or cause elusive bugs values[3] = 12;
- A common mistake: off-by-one error
$>$ Subscripts are between 0 and $n-1$, not 1 and $n$

```
int numbers[10];
    for (int count = 1; count <= 10; count++)
        numbers[count] = 0;
```


### 1.2 The Range-Based for Loop

- The range-based for loop is a loop that iterates once for each element in an array
- Each time the loop iterates, it copies an element from the array to a built-in variable, known as the range variable
- The range-based for loop automatically knows the number of elements in an array


## The Range-Based for Loop

- General format of the range-based for loop:

$$
\begin{aligned}
& \text { for (dataType rangeVariable : array) } \\
& \text { statement; }
\end{aligned}
$$

- dataType is the data type of the range variable.
- rangeVariable is the name of the range variable. This variable will receive the value of a different array element during each loop iteration.
- array is the name of an array.
- statement is a statement that executes during a loop iteration.


## Example

```
#include <iostream>
using namespace std;
int main() {
    // Define an array of integers
int numbers[] = {10, 20, 30, 40, 50};
    // Display the values in the array
for (int val : numbers)
    cout << val << endl;
}
return 0;
}
Output:

\subsection*{1.3 Processing Array Contents}
- Array elements can be treated as ordinary variables of the same type as the array
\(>\) Each element is a variable
\(>\) Processing an element is no different than processing other variables
- When using ++, -- operators, don't confuse the element with the subscript:
\[
\begin{aligned}
& \text { tests[i]++; // add } 1 \text { to tests[i] } \\
& \text { tests[i++]; // increment i, no } \\
& / / \text { effect on tests }
\end{aligned}
\]

\section*{Array Assignment}

\section*{To copy one array to another,}
- Don't try to assign one array to the other:
newTests = tests; // Won't work
- Instead, assign element-by-element:
\[
\begin{gathered}
\text { for }\left(i=0 ; i<A R R A Y \_S I Z E ; ~ i++\right) \\
\text { newTests[i] }=\text { tests[i]; }
\end{gathered}
\]

Note: Anytime the name of an array is used without brackets and a subscript, it is seen as the array's beginning memory address (not a variable).

\section*{In-class practice}
- Take 5 integers from user and store these numbers in an array
- Use a for loop to find the largest element of this array
- Display this element
- Test you code

Question: How to implement this practice using range-based for loop?

\subsection*{1.4 Arrays as Function Arguments}
- To pass an array to a function, use the array name:
\[
\begin{aligned}
& \text { int tests }[5]=\{79,82,91,77,84\} \text {; } \\
& \text { showScores(tests); }
\end{aligned}
\]
- To define a function that takes an array parameter, use empty [ ] for array argument:

> // function prototype
> void showScores(int []);
// function header No size declarator void showScores(int scores[f

Note: When an entire array is passed to a function, it is not passed by value, but passed by reference (only the starting memory address is passed).

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\section*{Arrays as Function Arguments}
- When passing an array to a function, it is common to pass array size so that function knows how many elements to process:

\author{
showScores(tests, ARRAY_SIZE); \# of elements
}
- Array size must also be reflected in prototype, header:
// function prototype
void showScores(int [], int);
// function header
void showScores(int scores[], int size)

\section*{Example}
```

\#include <iostream>
using namespace std;
void showValues(int [], int); // Function prototype
int main() {
const int ARRAY_SIZE = 8;
int numbers[ARRAY_SIZE] = {5, 10, 15, 20, 25, 30, 35, 40};
showValues(numbers, ARRAY_SIZE);
return 0;
}
void showValues(int nums[], int size) {
for (int index = 0; index < size; index++)
cout << nums[index] << " ";
cout << endl;
}
Output:
510152025303540

## In-class practice: Array Rotation

- Write a function Rotate that rotates an array of size $n$ by d elements to the left
- Use array as argument
- In the main function, call the function Rotate and show the rotated array
- Test your code

```
For example:
Input: [11 2 3 4 5 6 7], n = 7, d = 2
Output: [\begin{array}{llllllll}{3}&{4}&{5}&{6}&{7}&{1}&{2}\end{array}]
```

Reference code: ArrayRotation.cpp https://www.geeksforgeeks.org áandaydenotation/

### 1.5 Two-Dimensional Arrays

- A 2-D array is an array of 1-D arrays
- Use two size declarators in definition:
$>$ First declarator is number of rows; second is number of columns

```
const int ROWS = 4, COLS = 3;
int exams[ROWS][COLS];
```

columns

|  | exams[0][0] | exams [0][1] | exams [0][2] |
| :---: | :---: | :---: | :---: |
|  | exams[1][0] | exams[1][1] | exams[1][2] |
| W | exams[2][0] | exams[2][1] | exams [2][2] |
|  | exams[3][0] | exams[3][1] | exams[3][2] |

- Use two subscripts to access element:

```
exams[2][2] = 86;
```


## 2D Array Initialization

- Two-dimensional arrays are initialized row-by-row:

```
const int ROWS = 2, COLS = 2;
int exams[ROWS][COLS] = {{84, 78}, {92, 97}};
```

| 84 | 78 |
| :--- | :--- |
| 92 | 97 |

- Some array elements without initial values will be set to 0 or NULL

```
int exams[ROWS][COLS] = {{84}, {92, 97}};
```

exams[0][1] is automatically set to 0

## Passing Two-Dimensional Array to Function

- When a 2-D array is passed to a function, the parameter type must contain a size declarator for the columns
$>$ The size declarator for rows is optional (use empty [ ])

```
const int COLS = 2; Here COLS is a global constant
// Prototype
void getExams(int [][COLS], int);
```

// Header
void getExams(int exams[][COLS], int rows)

- Use array name as argument in function call: getExams (exams, 2);

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## Use Nested Loop to Step through 2D Array

- What is the output of the following program?

```
#include <iostream>
using namespace std;
```

int sum0fArray(int n[][2], int row) {

```
int sum0fArray(int n[][2], int row) {
    int total = 0;
    int total = 0;
    for (int i = 0; i < row; i++) {
    for (int i = 0; i < row; i++) {
        for (int j = 0; j< 2; j++) {
        for (int j = 0; j< 2; j++) {
        total += n[i][j];
        total += n[i][j];
        }
        }
    }
    }
    return total;
    return total;
}
```

}

```

Output:
The sum is: 29
\}
```

int main() {
int num[3][2]={{3, 4}, {9, 5}, {7, 1}};
cout <<" The sum is: " << sumOfArray (num, 3);
return 0;

```
https://www.programiz.com/cpp-programming/passing-arrays-function Salisbury

\section*{2. Searching and Sorting Arrays}
- 2.1 Array Search Algorithms
- 2.2 Array Sorting Algorithms

\subsection*{2.1 Array Search Algorithms}
- Search: locate an item in a list of data

- Two algorithms we will examine:
\(>\) Linear search
> Binary search
\({ }^{24}\) https://web.ics.purdue.edu/~cs154/lectures/lecture011.htm Salisbury

\section*{Linear search}
- Process
\(>\) Compare target \(\mathbf{x}\) with each element in an array in turn
> If \(\mathbf{x}\) matches with an element, return the index of this element
\(>\) If \(\mathbf{x}\) does not match with any elements, return -1


\section*{C++ implementation}
```

int linearSearch(int arr[], int size, int value)
{
int index = 0; // Search index
int position = -1; // Location of the value
bool found = false; // Search flag
while (index < size \&\& !found)
{
if (arr[index] == value) // Value is found
{
found = true; // Set the flag
position = index; // Record the location
}
index++;
}
return position;
}

```

\section*{Linear Search - Tradeoffs}
- Benefits:
\(>\) Easy algorithm to understand
\(>\) Array can be in any order
- Disadvantages:
\(>\) Inefficient (slow): for array of N elements, examines \(\mathrm{N} / 2\) elements on average for value in array, N elements for value not in array

\section*{Binary search (Example)}

Value \(=6\)
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|}
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\
\hline \(\mathbf{1}\) & \(\mathbf{2}\) & \(\mathbf{3}\) & \(\mathbf{4}\) & \(\mathbf{5}\) & \(\mathbf{6}\) & \(\mathbf{7}\) & \(\mathbf{8}\) & \(\mathbf{9}\) & \(\mathbf{1 0}\) \\
\hline
\end{tabular}
\begin{tabular}{l|l|l|l|l|l|l|l|l|l|l|}
\hline \(6>5\) \\
Take \(2^{\text {nd }}\) half & \(\mathbf{1}\) & \(\mathbf{2}\) & \(\mathbf{3}\) & \(\mathbf{4}\) & 5 & \(\mathbf{6}\) & \(\mathbf{7}\) & \(\mathbf{8}\) & \(\mathbf{9}\) & \(\mathbf{1 0}\) \\
\hline
\end{tabular}
\(6<8\)
Take \(1^{\text {st }}\) half
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|}
\hline & 0 & 1 & 2 & 3 & 4 & \(\mathrm{L}=5\) & 6 & M & 8 & H = 9 \\
\hline \begin{tabular}{l}
\[
6<8
\] \\
Take \(1^{\text {st }}\) half
\end{tabular} & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \\
\hline
\end{tabular}

Found 6
Return index 5
\[
\begin{array}{llllllllll}
\mathrm{L}=0 & 1 & 2 & 3 & \mathbf{M}=\mathbf{4} & 5 & 6 & 7 & 8 & \mathbf{H}=9
\end{array}
\]
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|}
\hline & 0 & 1 & 2 & 3 & 4 & \[
\begin{aligned}
& \mathrm{L}= \\
& \mathrm{M}=
\end{aligned}
\] & H & 7 & 8 & 9 \\
\hline Found 6 Return index 5 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \\
\hline
\end{tabular}

\section*{Process of binary search}
- Step 1: find the middle element, middle
- Step 2: compare middle with value
\(>\) If value < middle, drop the second half
\(>\) If value > middle, drop the first half
\(>\) If value \(==\) middle, the search is finished
- Repeat above steps. If no element left, value is not in the array

\section*{C++ implementation}
```

int binarySearch(int array[], int size, int value)
{
int first = 0, // First array element
last = size - 1,
middle,
position = -1;
bool found = false;
while (!found \&\& first <= last)
{
middle = (first + last) / 2; // Middle point
if (array[middle] == value) // If value = middle
{
found = true;
position = middle;
}
else if (array[middle] > value) // If value < middle
last = middle - 1; // Search lower half
else
first = middle + 1; // If value > middle
}
// Search upper half
return position;
}

## Binary Search - Tradeoffs

- Benefits:
$>$ Much more efficient than linear search. For array of $N$ elements, performs at most $\log _{2} N$ comparisons

- Disadvantages:
$>$ Requires that array elements be sorted


## Linear search VS binary search

## Linear search

+ No need to sort elements

Only equality comparisons
Sequential access to the data

- Search is inefficient (slow)


## Binary search

- Need to sort elements first

Equality \& ordering comparisons
Random access to the data

+ Search is efficient (fast)


## In-class practice

- Search Insert Position
- Given a sorted array in ascending order and a target value
- Use binary search algorithm to return the index if the target is found. If not, return the index where it would be if it is inserted in order
- You may assume no duplicates in the array


## Example :

Input: [1,3,5,6], 5
Output: 2
Input: [1,3,5,6], 2
Output: 1

### 2.2 Array Sorting Algorithms

- Sort: arrange values into an order
$>$ Alphabetical
$>$ Ascending numeric
$>$ Descending numeric

- Two algorithms considered here:
$>$ Bubble sort
$>$ Selection sort


## Bubble Sort (Example)

- Sort an array in ascending order


## First pass:



| 2 | 3 | 7 | 8 | 1 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- |

## Second pass:



| 2 | 3 | 7 | 1 | 8 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- |

The $2^{\text {nd }}$ largest value 8 is in the correct position

Note: the $2^{\text {nd }}$ pass will not involve the last element bcz 9 is the largest

## Bubble Sort

Third pass:


| 2 | 3 | 1 | 7 | 8 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- |

The $3^{\text {rd }}$ largest value 7 is in the correct position

Note: the $3^{\text {rd }}$ pass will not involve the last two elements bcz they are sorted

## Fourth pass:

$2<3$, no swap

| 2 | 3 | 1 | 7 | 8 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- |



| 2 | 1 | 3 | 7 | 8 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- |

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## Bubble Sort

## Fifth pass:



| 1 | 2 | 3 | 7 | 8 | 9 |
| :--- | :--- | :--- | :--- | :--- | :--- |

- There are $(\mathrm{n}-1)$ passes. n is the number of elements in the array

| Pass | 1 | 2 | $\ldots$ | $n-2$ | $n-1$ |
| :--- | :---: | :---: | :---: | :---: | :---: |
| \# of compares | $\mathrm{n}-1$ | $\mathrm{n}-2$ | $\ldots$ | 2 | 1 |

In total $(\mathrm{n}-1)+(\mathrm{n}-2)+\ldots+(2)+(1)=\mathrm{n}(\mathrm{n}-1) / 2$ comparisons.

## Bubble Sort

## Process:

$>$ Compare $1^{\text {st }}$ and $2^{\text {nd }}$ elements

- If out of order, exchange them to put in order
$>$ Move down one element, compare $2^{\text {nd }}$ and $3^{\text {rd }}$ elements, exchange if necessary. Continue until end of array
$>$ Pass through array (one element less) again, exchanging as necessary
>Repeat until the last pass


## C++ Implementation

```
void bubbleSort(int array[], int size)
    int maxElement; Hold the subscript of the last element to be compared
    int index;
                                    Used as an array subscript in one of the loops
    for (maxElement = size - 1; maxElement > 0; maxElement--) {
        for (index = 0; index < maxElement; index++) {
        if (array[index] > array[index + 1]) {
                        swap(array[index], array[index + 1]);
        }
        }
    }
}
void swap(int &a, int &b)fmamomen, Reference parameters
    int temp = a;
    a = b;
    b = temp;
}
```


## Bubble Sort - Tradeoffs

- Benefit:
>Easy to understand and implement
- Disadvantage:
>Inefficient: slow for large arrays
- Too much unnecessary swaps

Question: How many swaps for bubble sort in the worst case?
$n(n-1) / 2$ when the array is reversely sorted

## Selection Sort

- Concept for sort in ascending order:
$>$ Locate smallest element in array.
Exchange it with element in position 0
$>$ Locate next smallest element in array.
Exchange it with element in position 1.
>Continue until all elements are arranged in order


## Selection Sort - Example

Array numlist contains:


1. Smallest element is 2 . Exchange 2 with element in $1^{\text {st }}$ position in array:


## Example (Continued)

2. Next smallest element is 3 . Exchange 3 with element in $2^{\text {nd }}$ position in array:

3. Next smallest element is 11. Exchange 11 with element in $3^{\text {rd }}$ position in array:


## C++ Implementation

```
void selectionSort(int array[], int size) {
    int minIndex, minValue;
    for (int start = 0; start < (size - 1); start++) {
    minIndex = start;
    minValue = array[start];
    for (int index = start + 1; index < size; index++) {
        if (array[index] < minValue) {
            minValue = array[index];
            minIndex = index;
        }
    }
    swap(array[minIndex], array[start]);
    }
```


## Selection Sort - Tradeoffs

- Benefit:
$>$ More efficient than Bubble Sort, since fewer exchanges/swaps
- Disadvantage:
$>$ May not be as easy as Bubble Sort to understand

Question: How many comparisons for selection sort? In total $(n-1)+(n-2)+\ldots+(2)+(1)=n(n-1) / 2$ comparisons.

Question: How many swaps for selection sort in the worst case?
( $\mathrm{n}-1$ ) when the array is reversely sorted

## In-class practice

- For an array with n elements, the bubble sort needs $n-1$ passes. However, if the array elements are in order in the midway, there is no need to execute the subsequent passes.
- Write code to implement the above optimized bubble sort algorithm.
- Test you code.


## Recursion Function

- A recursive function is one that calls itself

```
void countDown(int num) {
    if (num == 0) // stop condition
        cout << "Blastoff!";
        else{
        cout << num << "...\n";
        countDown (num-1);
                            // recursive call
}
```

- Assume the input argument is 2 :
$>$ countDown (2) outputs 2 . . ., then it calls countDown (1)
$>$ countDown (1) outputs 1 . .., then it calls countDown (0)
$>$ countDown (0) outputs Blastoff!, then returns to countDown (1)
$>$ countDown (1) returns to countDown (2)
$>$ countDown (2) returns to the calling function


## What Happens When Called?



## Solving Problems with Recursion

- Two important steps:
$>$ Define the recursive function
$>$ Define the stop condition
- Example: factorial calculation

$$
\begin{aligned}
& n!=n *(n-1) *(n-2) * \ldots * 3 * 2 * 1 \text { if } n>0 \\
& n!=1 \text { if } n=0
\end{aligned}
$$

$>$ Define the recursive function:

$$
n!=n *(n-1)!
$$

$>$ Define the stop condition:
$0!=1$ (base case)

## Recursive Factorial Function

```
#include <iostream>
using namespace std;
int factorial(int); // Function prototype
int main(){
    int number;
    cout << "Enter an integer value to display its factorial: ";
    cin >> number;
    cout << "The factorial of " << number << " is " << factorial(number);
    return 0;
}
int factorial(int n){
    if (n == 0)
        return 1; // Base case
    else
        return n * factorial(n-1); // Recursive case
}
```

Enter an integer value to display its factorial: 5
The factorial of 5 is 120

## In-class practice

- The Fibonacci numbers are the numbers in the following integer sequence
$0,1,1,2,3,5,8,13,21,34,55,89,144, \ldots$
- In mathematical terms, the sequence Fn of Fibonacci numbers is defined as
$\mathrm{Fn}=\mathrm{Fn}-1+\mathrm{Fn}-2$
where:
$\mathrm{FO}=0$ and $\mathrm{F} 1=1$
- Use recursive function to calculate and display the first 10 Fibonacci numbers
- Test your code


## Application: The Towers of Hanoi

- The game uses three pegs and a set of discs, stacked on one of the pegs
- The object of this game is to move the discs from the first peg to the third peg
- Here are the rules:
$>$ Only one disc may be moved at a time
$>$ A disc cannot be placed on top of a smaller disc
$>$ All discs must be stored on a peg except while being moved



## Moving Three Discs



Original setup.


Second move: Move disc 2 to peg 2.


Fourth move: Move disc 3 to peg 3.


Sixth move: Move disc 2 to peg 3 .


First move: Move disc 1 to peg 3.


Third move: Move disc 1 to peg 2.


Fifth move: Move disc 1 to peg 1.


Seventh move: Move disc 1 to peg 3.

## The Towers of Hanoi

- Algorithm
$>$ To move $n$ discs from peg $A$ to peg $C$, using peg $B$ as a temporary peg: If $n>0$ Then
Move $n$ - 1 discs from peg $A$ to peg B, using peg $C$ as a temporary peg.

Move the remaining disc from the peg $A$ to peg $C$.
Move n - 1 discs from peg B to peg C, using peg $A$ as a temporary peg.

End If

## The Towers of Hanoi

- C++ Implementation
$>$ Refer to "Pr20-10.cpp"


## Reference

- The teaching materials of this course refer to:
$>$ Professor Xiaohong (Sophie) Wang. COSC 120 teaching materials
- Salisbury University
$>$ Textbook:
- Starting Out with C++: From Control Structures through Objects, by Tony Gaddis, Pearson (9th Edition)
- Instructor materials of the above textbook (All rights reserved)

